

CS61C : Machine Structures

Lecture 2 – Number Representation



2004-09-01

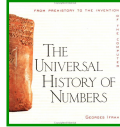
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Great book ⇒
The Universal History
of Numbers

by Georges Ifrah

Happy September 1st!



Decimal Numbers: Base 10

Digits: 0, 1, 2, 3, 4, 5, 6, 7, 8, 9

Example:

3271 =

$$(3 \times 10^3) + (2 \times 10^2) + (7 \times 10^1) + (1 \times 10^0)$$



Numbers: positional notation

• Number Base B ⇒ B symbols per digit:

- Base 10 (Decimal): 0, 1, 2, 3, 4, 5, 6, 7, 8, 9
- Base 2 (Binary): 0, 1

• Number representation:

- $d_{31}d_{30} \dots d_1d_0$ is a 32 digit number
- value = $d_{31} \times B^{31} + d_{30} \times B^{30} + \dots + d_1 \times B^1 + d_0 \times B^0$

• Binary: 0,1 (In binary digits called “bits”)

• $0b11010 = 1 \times 2^4 + 1 \times 2^3 + 0 \times 2^2 + 1 \times 2^1 + 0 \times 2^0$
 $= 16 + 8 + 2$
 #s often written = 26

- Here 5 digit binary # turns into a 2 digit decimal #
- Can we find a base that converts to binary easily?



Hexadecimal Numbers: Base 16

• Hexadecimal:

0, 1, 2, 3, 4, 5, 6, 7, 8, 9, A, B, C, D, E, F

- Normal digits + 6 more from the alphabet

- In C, written as 0x... (e.g., 0xFAB5)

• Conversion: Binary ⇔ Hex

- 1 hex digit represents 16 decimal values
- 4 binary digits represent 16 decimal values
- ⇒ 1 hex digit replaces 4 binary digits

• One hex digit is a “nibble”. Two is a “byte”

• Example:

- 1010 1100 0011 (binary) = 0x _____ ?



Decimal vs. Hexadecimal vs. Binary

Examples:

1010 1100 0011 (binary)
= 0xAC3

10111 (binary)
= 0001 0111 (binary)
= 0x17

0x3F9
= 11 1111 1001 (binary)

How do we convert between
hex and Decimal?

MEMORIZE!

00	0	0000
01	1	0001
02	2	0010
03	3	0011
04	4	0100
05	5	0101
06	6	0110
07	7	0111
08	8	1000
09	9	1001
10	A	1010
11	B	1011
12	C	1100
13	D	1101
14	E	1110
15	F	1111



What to do with representations of numbers?

• Just what we do with numbers!

- Add them $\begin{array}{r} 1 \quad 1 \\ \end{array}$
- Subtract them $\begin{array}{r} 1 \quad 0 \quad 1 \quad 0 \\ \end{array}$
- Multiply them $\begin{array}{r} + \quad 0 \quad 1 \quad 1 \quad 1 \\ \end{array}$
- Divide them $\begin{array}{r} \quad \quad \quad \quad 1 \\ \hline \quad \quad \quad \quad 1 \end{array}$
- Compare them $\begin{array}{r} 1 \quad 0 \quad 0 \quad 0 \quad 1 \\ \end{array}$

• Example: 10 + 7 = 17

- ...so simple to add in binary that we can build circuits to do it!
- subtraction just as you would in decimal
- Comparison: How do you tell if X > Y ?



Which base do we use?


- **Decimal:** great for humans, especially when doing arithmetic
- **Hex:** if human looking at long strings of binary numbers, its much easier to convert to hex and look 4 bits/symbol
 - Terrible for arithmetic on paper
- **Binary:** what computers use; you will learn how computers do +, -, *, /
 - To a computer, numbers always binary
 - Regardless of how number is written:
 $32_{ten} == 32_{10} == 0x20 == 100000_2 == 0b100000$
 - Use subscripts "ten", "hex", "two" in book, slides when might be confusing



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BIG IDEA: Bits can represent anything!!

- **Characters?**
 - 26 letters \Rightarrow 5 bits ($2^5 = 32$)
 - upper/lower case + punctuation \Rightarrow 7 bits (in 8) ("ASCII")
 - standard code to cover all the world's languages \Rightarrow 8,16,32 bits ("Unicode")
www.unicode.com 
- **Logical values?**
 - 0 \Rightarrow False, 1 \Rightarrow True
- **colors ? Ex:** Red (00) Green (01) Blue (11)
- **locations / addresses? commands?**
- **MEMORIZE: N bits \Rightarrow at most 2^N things**



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How to Represent Negative Numbers?

- So far, **unsigned numbers**
- **Obvious solution:** define leftmost bit to be sign!
 - 0 \Rightarrow +, 1 \Rightarrow -
 - Rest of bits can be numerical value of number
- Representation called **sign and magnitude**
- MIPS uses 32-bit integers. $+1_{ten}$ would be:
0000 0000 0000 0000 0000 0000 0001
- And -1_{ten} in sign and magnitude would be:
1000 0000 0000 0000 0000 0000 0001



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Shortcomings of sign and magnitude?

- **Arithmetic circuit complicated**
 - Special steps depending whether signs are the same or not
- **Also, two zeros**
 - $0x00000000 = +0_{ten}$
 - $0x80000000 = -0_{ten}$
 - What would two 0s mean for programming?
- **Therefore sign and magnitude abandoned**



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Administrivia

- **Firmware (n) :** a set of computer instructions used so frequently that it is stored on a memory chip in a computer rather than being part of a program. [MS Word Dictionary]
- Look at class website often!
- Reading for next Wed: K&R Ch1-4
 - 1st Reading quiz due Wed [if not up by 6pm the day before, it's not due & will be postponed]
- Labs need to be finished in **your lab session** (& only your lab) unless extended by your TA
 - +1 bonus pt if you finish your lab in the first hour!
- **Homework #1** up now, due Wed @ 11:59pm



Homework #2 up soon, due following Wed

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Another try: complement the bits

- **Example:** $7_{10} = 00111_2$ $-7_{10} = 11000_2$
- Called **One's Complement**
- **Note:** positive numbers have leading 0s, negative numbers have leading 1s.
$$\begin{array}{ccccccc} & 00000 & 00001 & \dots & 01111 & & \\ \leftarrow & & & & & & \rightarrow \\ & 10000 & \dots & 11110 & 11111 & & \end{array}$$
- What is -00000 ? Answer: 11111
- How many positive numbers in N bits?
- How many negative ones?



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Shortcomings of One's complement?

- Arithmetic still a somewhat complicated.
- Still two zeros
 - $0x00000000 = +0_{ten}$
 - $0xFFFFFFF = -0_{ten}$
- Although used for awhile on some computer products, one's complement was eventually abandoned because another solution was better.



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Standard Negative Number Representation

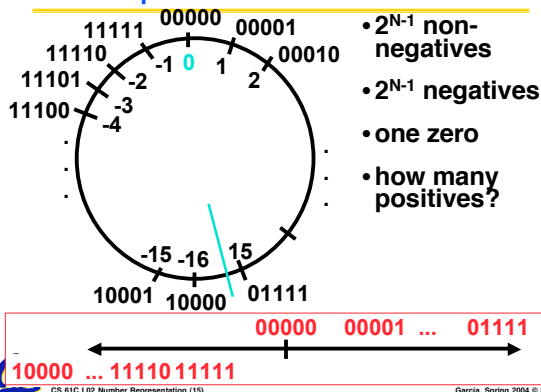
- What is result for unsigned numbers if tried to subtract large number from a small one?
 - Would try to borrow from string of leading 0s, so result would have a string of leading 1s
 - $3 - 4 \Rightarrow 00...0011 - 00...0100 = 11...1111$
 - With no obvious better alternative, pick representation that made the hardware simple
 - As with sign and magnitude, leading 0s \Rightarrow positive, leading 1s \Rightarrow negative
 - $000000...xxx$ is ≥ 0 , $111111...xxx$ is < 0
 - except $1...1111$ is -1 , not -0 (as in sign & mag.)
- This representation is **Two's Complement**



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2's Complement Number "line": N = 5



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Two's Complement for N=32

0000 ... 0000 0000 0000 0000	$_{two} =$	0	$_{ten}$
0000 ... 0000 0000 0000 0001	$_{two} =$	1	$_{ten}$
0000 ... 0000 0000 0000 0010	$_{two} =$	2	$_{ten}$
0111 ... 1111 1111 1111 1101	$_{two} =$	2,147,483,645	$_{ten}$
0111 ... 1111 1111 1111 1110	$_{two} =$	2,147,483,646	$_{ten}$
0111 ... 1111 1111 1111 1111	$_{two} =$	2,147,483,647	$_{ten}$
1000 ... 0000 0000 0000 0000	$_{two} =$	-2,147,483,648	$_{ten}$
1000 ... 0000 0000 0000 0001	$_{two} =$	-2,147,483,647	$_{ten}$
1000 ... 0000 0000 0000 0010	$_{two} =$	-2,147,483,646	$_{ten}$
1111 ... 1111 1111 1111 1101	$_{two} =$	-3	$_{ten}$
1111 ... 1111 1111 1111 1110	$_{two} =$	-2	$_{ten}$
1111 ... 1111 1111 1111 1111	$_{two} =$	-1	$_{ten}$

- One zero; 1st bit called **sign bit**
- 1 "extra" negative: no positive $2,147,483,648_{ten}$



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Two's Complement Formula

- Can represent positive **and negative** numbers in terms of the bit value times a power of 2:

$$d_{31} \times (-2^{31}) + d_{30} \times 2^{30} + \dots + d_2 \times 2^2 + d_1 \times 2^1 + d_0 \times 2^0$$
- Example: 1101_{two}

$$= 1x(-2^3) + 1x2^2 + 0x2^1 + 1x2^0$$

$$= -2^3 + 2^2 + 0 + 2^0$$

$$= -8 + 4 + 0 + 1$$

$$= -8 + 5$$

$$= -3_{ten}$$



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Two's Complement shortcut: Negation

- Change every 0 to 1 and 1 to 0 (invert or complement), then add 1 to the result
- Proof: Sum of number and its (one's) complement must be $111...111_{two}$
 - However, $111...111_{two} = -1_{ten}$
 - Let $x' \Rightarrow$ one's complement representation of x
 - Then $x + x' = -1 \Rightarrow x + x' + 1 = 0 \Rightarrow x' + 1 = -x$

• Example: -3 to +3 to -3

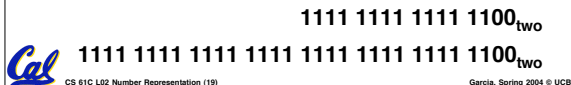
x:	1111 1111 1111 1111 1111 1111 1111 1101	$_{two}$
x':	0000 0000 0000 0000 0000 0000 0000 0010	$_{two}$
+1:	0000 0000 0000 0000 0000 0000 0000 0011	$_{two}$
():	1111 1111 1111 1111 1111 1111 1111 1100	$_{two}$
+1:	1111 1111 1111 1111 1111 1111 1111 1101	$_{two}$



You should be able to do this in your head... © UCB

Two's comp. shortcut: Sign extension

- Convert 2's complement number rep. using n bits to more than n bits
- Simply **replicate** the most significant bit (sign bit) of smaller to fill new bits
 - 2's comp. positive number has infinite 0s
 - 2's comp. negative number has infinite 1s
 - Binary representation hides leading bits; sign extension restores some of them
 - 16-bit -4_{ten} to 32-bit:

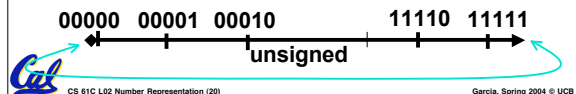


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What if too big?

- Binary bit patterns above are simply **representatives** of numbers. Strictly speaking they are called "numerals".
- Numbers really have an ∞ number of digits
 - with almost all being same (00...0 or 11...1) except for a few of the rightmost digits
 - Just don't normally show leading digits
- If result of add (or -, *, /) cannot be represented by these rightmost HW bits, **overflow** is said to have occurred.

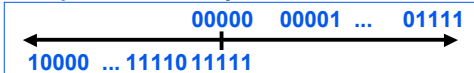


CS 61C L02 Number Representation (20)

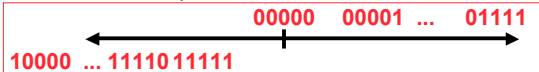
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And in Conclusion...

- We represent "things" in computers as particular bit patterns: $N \text{ bits} \Rightarrow 2^N$
- Decimal for human calculations, binary for computers, hex to write binary more easily
- 1's complement - mostly abandoned



- 2's complement universal in computing: cannot avoid, so learn



Overflow: numbers ∞ ; computers finite, errors!



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Bonus Slides

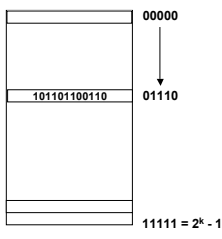
- Peer instruction let's us skip example slides since you are expected to read book and lecture notes beforehand, but we include them for your review
- Slides shown in logical sequence order



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BONUS: Numbers represented in memory



- Memory is a place to store bits
- A *word* is a fixed number of bits (eg, 32) at an address
- *Addresses* are naturally represented as unsigned numbers in C



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BONUS: Signed vs. Unsigned Variables

- Java just declares integers `int`
 - Uses two's complement
- C has declaration `int` also
 - Declares variable as a signed integer
 - Uses two's complement
- Also, C declaration `unsigned int`
 - Declares a unsigned integer
 - Treats 32-bit number as unsigned integer, so most significant bit **is part of the number**, not a sign bit



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