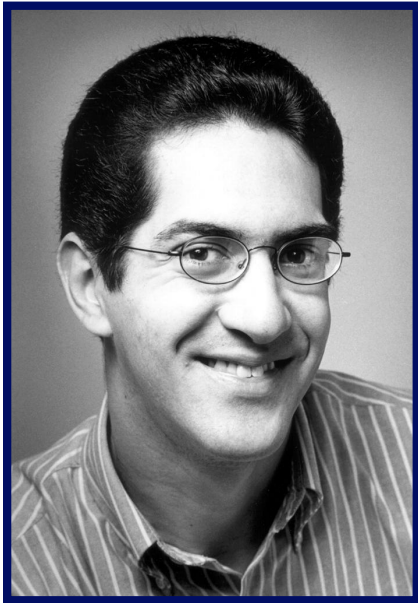


inst.eecs.berkeley.edu/~cs61c

# CS61C : Machine Structures

## Lecture 2 – Number Representation

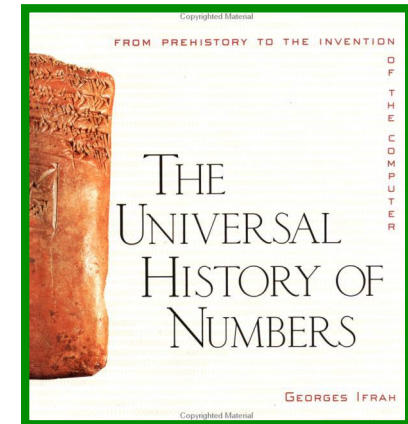


2004-09-01

Lecturer PSOE Dan Garcia

[www.cs.berkeley.edu/~ddgarcia](http://www.cs.berkeley.edu/~ddgarcia)

Great book ⇒  
**The Universal History  
of Numbers**  
by **Georges Ifrah**



Happy September 1<sup>st</sup>!



# Decimal Numbers: Base 10

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**Digits: 0, 1, 2, 3, 4, 5, 6, 7, 8, 9**

**Example:**

**3271 =**

$$(3 \times 10^3) + (2 \times 10^2) + (7 \times 10^1) + (1 \times 10^0)$$



# Numbers: positional notation

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- Number Base  $B \Rightarrow B$  symbols per digit:

- Base 10 (Decimal): 0, 1, 2, 3, 4, 5, 6, 7, 8, 9
- Base 2 (Binary): 0, 1

- Number representation:

- $d_{31}d_{30} \dots d_1d_0$  is a 32 digit number

- value =  $d_{31} \times B^{31} + d_{30} \times B^{30} + \dots + d_1 \times B^1 + d_0 \times B^0$

- Binary: 0,1 (In binary digits called “bits”)



- $0b11010 = 1 \times 2^4 + 1 \times 2^3 + 0 \times 2^2 + 1 \times 2^1 + 0 \times 2^0$   
 $= 16 + 8 + 2$   
#s often written = 26

- Here 5 digit binary # turns into a 2 digit decimal #

- Can we find a base that converts to binary easily?



# Hexadecimal Numbers: Base 16

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- **Hexadecimal:**  
0, 1, 2, 3, 4, 5, 6, 7, 8, 9, A, B, C, D, E, F
  - Normal digits + 6 more from the alphabet
  - In C, written as **0x...** (e.g., 0xFAB5)
- **Conversion: Binary  $\Leftrightarrow$  Hex**
  - 1 hex digit represents 16 decimal values
  - 4 binary digits represent 16 decimal values
  - $\Rightarrow$  1 hex digit replaces 4 binary digits
- One hex digit is a “**nibble**”. Two is a “**byte**”
- **Example:**
  - 1010 1100 0011 (binary) = **0x\_\_\_\_\_** ?



# Decimal vs. Hexadecimal vs. Binary

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**Examples:**

**1010 1100 0011 (binary)  
= 0xAC3**

**10111 (binary)  
= 0001 0111 (binary)  
= 0x17**

**0x3F9  
= 11 1111 1001 (binary)**

*How do we convert between  
hex and Decimal?*

00	0	0000
01	1	0001
02	2	0010
03	3	0011
04	4	0100
05	5	0101
06	6	0110
07	7	0111
08	8	1000
09	9	1001
10	A	1010
11	B	1011
12	C	1100
13	D	1101
14	E	1110
15	F	1111

**MEMORIZE!**



# What to do with representations of numbers?

- **Just what we do with numbers!**

- Add them
- Subtract them
- Multiply them
- Divide them
- Compare them

$$\begin{array}{r} \phantom{+} 1010 \\ + 0111 \\ \hline 10001 \end{array}$$

- **Example:  $10 + 7 = 17$**

- ...so simple to add in binary that we can build circuits to do it!
- subtraction just as you would in decimal
- Comparison: How do you tell if  $X > Y$  ?



# Which base do we use?

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- **Decimal: great for humans, especially when doing arithmetic**
- **Hex: if human looking at long strings of binary numbers, its much easier to convert to hex and look 4 bits/symbol**
  - Terrible for arithmetic on paper
- **Binary: what computers use; you will learn how computers do +, -, \*, /**
  - To a computer, numbers always binary
  - Regardless of how number is written:  
 $32_{\text{ten}} == 32_{10} == 0x20 == 100000_2 == 0b100000$
  - Use subscripts “ten”, “hex”, “two” in book, slides when might be confusing



# BIG IDEA: Bits can represent anything!!

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- **Characters?**

- 26 letters  $\Rightarrow$  5 bits ( $2^5 = 32$ )
- upper/lower case + punctuation  $\Rightarrow$  7 bits (in 8) (“ASCII”)
- standard code to cover all the world’s languages  $\Rightarrow$  8,16,32 bits (“Unicode”) [www.unicode.com](http://www.unicode.com)



- **Logical values?**

- 0  $\Rightarrow$  False, 1  $\Rightarrow$  True

- **colors ? Ex:** Red (00) Green (01) Blue (11)

- **locations / addresses? commands?**

- **MEMORIZE: N bits  $\Rightarrow$  at most  $2^N$  things**





# How to Represent Negative Numbers?

- So far, **unsigned numbers**
- Obvious solution: define leftmost bit to be sign!
  - $0 \Rightarrow +$ ,  $1 \Rightarrow -$
  - Rest of bits can be numerical value of number
- Representation called **sign and magnitude**
- MIPS uses 32-bit integers.  $+1_{\text{ten}}$  would be:  
**0000 0000 0000 0000 0000 0000 0000 0001**
- And  $-1_{\text{ten}}$  in sign and magnitude would be:  
**1000 0000 0000 0000 0000 0000 0000 0001**



# Shortcomings of sign and magnitude?

---

- **Arithmetic circuit complicated**
  - Special steps depending whether signs are the same or not
- **Also, two zeros**
  - $0x00000000 = +0_{\text{ten}}$
  - $0x80000000 = -0_{\text{ten}}$
  - What would two 0s mean for programming?
- **Therefore sign and magnitude abandoned**



# Administrivia

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- **Firmware (n) : a set of computer instructions used so frequently that it is stored on a memory chip in a computer rather than being part of a program. [MS Word Dictionary]**
- **Look at class website often!**
- **Reading for next Wed: K&R Ch1-4**
  - **1<sup>st</sup> Reading quiz due Wed [if not up by 6pm the day before, it's not due & will be postponed]**
- **Labs need to be finished in **your lab session** (& only your lab) unless extended by your TA**
  - **+1 bonus pt if you finish your lab in the first hour!**
- **Homework #1 up now, due Wed @ 11:59pm**

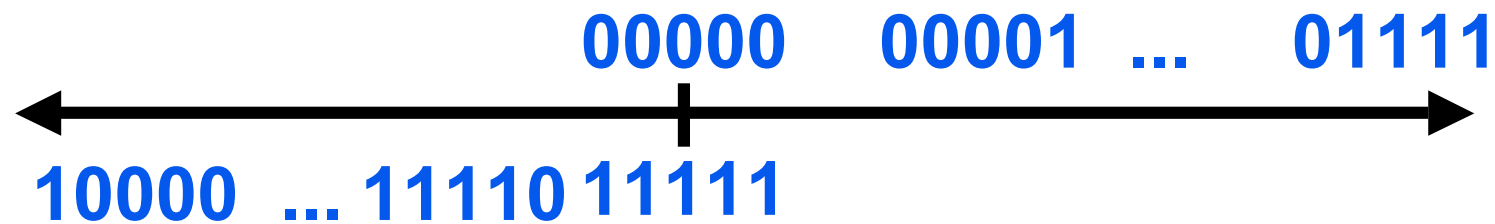


• **Homework #2 up soon, due following Wed**

## Another try: complement the bits

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- Example:  $7_{10} = 00111_2$   $-7_{10} = 11000_2$
- Called One's Complement
- Note: positive numbers have leading 0s, negative numbers have leading 1s.



- What is  $-00000$  ? Answer:  $11111$
- How many positive numbers in N bits?
- How many negative ones?



# Shortcomings of One's complement?

---

- **Arithmetic still a somewhat complicated.**
- **Still two zeros**
  - $0x00000000 = +0_{\text{ten}}$
  - $0xFFFFFFFF = -0_{\text{ten}}$
- **Although used for awhile on some computer products, one's complement was eventually abandoned because another solution was better.**

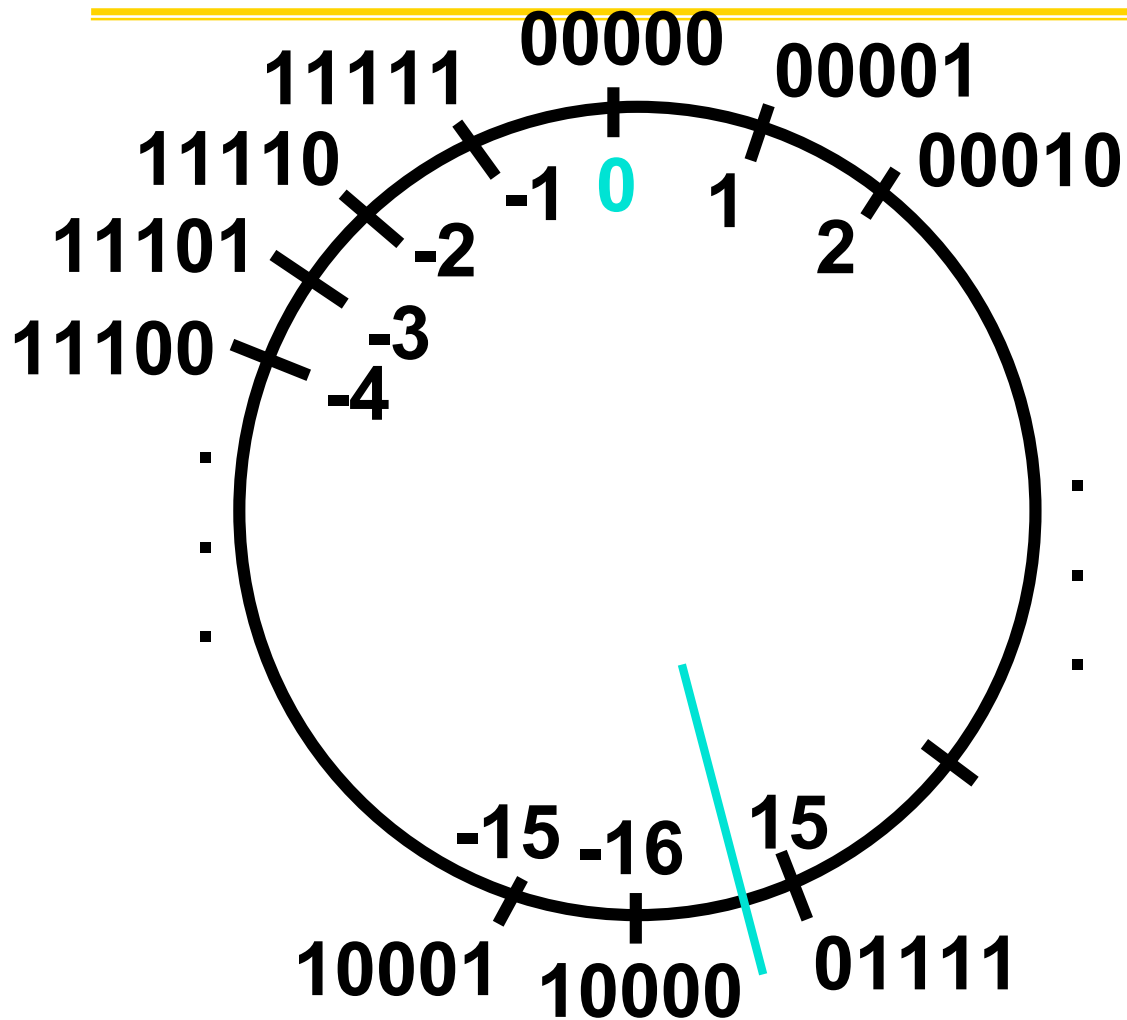


# Standard Negative Number Representation

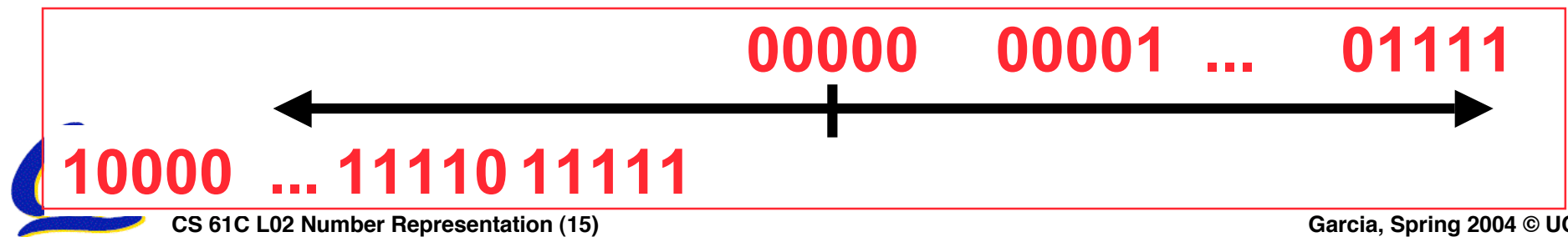
- What is result for unsigned numbers if tried to subtract large number from a small one?
  - Would try to borrow from string of leading 0s, so result would have a string of leading 1s
    - $3 - 4 \Rightarrow 00\dots0011 - 00\dots0100 = 11\dots1111$
  - With no obvious better alternative, pick representation that made the hardware simple
  - As with sign and magnitude, leading 0s  $\Rightarrow$  positive, leading 1s  $\Rightarrow$  negative
    - $000000\dots xxx$  is  $\geq 0$ ,  $111111\dots xxx$  is  $< 0$
    - except  $1\dots1111$  is -1, not -0 (as in sign & mag.)
- This representation is Two's Complement



## 2's Complement Number "line": N = 5



- $2^{N-1}$  non-negatives
- $2^{N-1}$  negatives
- one zero
- how many positives?



# Two's Complement for N=32

0000 ... 0000 0000 0000 0000	$\text{two} =$	$0_{\text{ten}}$
0000 ... 0000 0000 0000 0001	$\text{two} =$	$1_{\text{ten}}$
0000 ... 0000 0000 0000 0010	$\text{two} =$	$2_{\text{ten}}$
$\ddots$ 0111 ... 1111 1111 1111 1101	$\text{two} =$	$2,147,483,645_{\text{ten}}$
0111 ... 1111 1111 1111 1110	$\text{two} =$	$2,147,483,646_{\text{ten}}$
0111 ... 1111 1111 1111 1111	$\text{two} =$	$2,147,483,647_{\text{ten}}$
1000 ... 0000 0000 0000 0000	$\text{two} =$	$-2,147,483,648_{\text{ten}}$
1000 ... 0000 0000 0000 0001	$\text{two} =$	$-2,147,483,647_{\text{ten}}$
1000 ... 0000 0000 0000 0010	$\text{two} =$	$-2,147,483,646_{\text{ten}}$
$\ddots$ 1111 ... 1111 1111 1111 1101	$\text{two} =$	$-3_{\text{ten}}$
1111 ... 1111 1111 1111 1110	$\text{two} =$	$-2_{\text{ten}}$
1111 ... 1111 1111 1111 1111	$\text{two} =$	$-1_{\text{ten}}$

- One zero; 1st bit called **sign bit**
- 1 “extra” negative: no positive  $2,147,483,648_{\text{ten}}$





# Two's Complement Formula

- Can represent positive and negative numbers in terms of the bit value times a power of 2:

$$d_{31} \times \text{-(2}^{31}\text{)} + d_{30} \times 2^{30} + \dots + d_2 \times 2^2 + d_1 \times 2^1 + d_0 \times 2^0$$

- Example:  $1101_{\text{two}}$

$$= 1 \times \text{-(2}^3\text{)} + 1 \times 2^2 + 0 \times 2^1 + 1 \times 2^0$$

$$= \text{-2}^3 + 2^2 + 0 + 2^0$$

$$= \text{-8} + 4 + 0 + 1$$

$$= \text{-8} + 5$$

$$= \text{-3}_{\text{ten}}$$



## Two's Complement shortcut: Negation

- Change every 0 to 1 and 1 to 0 (invert or complement), then add 1 to the result

- Proof: Sum of number and its (one's) complement must be  $111\dots111_{\text{two}}$

However,  $111\dots111_{\text{two}} = -1_{\text{ten}}$

Let  $x' \Rightarrow$  one's complement representation of  $x$

Then  $x + x' = -1 \Rightarrow x + x' + 1 = 0 \Rightarrow x' + 1 = -x$

- Example: -3 to +3 to -3

$x$	:	1111	1111	1111	1111	1111	1111	1111	1101	$_{\text{two}}$
$x'$	:	0000	0000	0000	0000	0000	0000	0000	0010	$_{\text{two}}$
$+1$	:	0000	0000	0000	0000	0000	0000	0000	0011	$_{\text{two}}$
$( )'$	:	1111	1111	1111	1111	1111	1111	1111	1100	$_{\text{two}}$
$+1$	:	1111	1111	1111	1111	1111	1111	1111	1101	$_{\text{two}}$



You should be able to do this in your head...

# Two's comp. shortcut: Sign extension

- Convert 2's complement number rep. using  $n$  bits to more than  $n$  bits
- Simply **replicate** the most significant bit (sign bit) of smaller to fill new bits
  - 2's comp. positive number has infinite 0s
  - 2's comp. negative number has infinite 1s
  - Binary representation hides leading bits; sign extension restores some of them
  - 16-bit  $-4_{\text{ten}}$  to 32-bit:

1111 1111 1111 1100<sub>two</sub>

1111 1111 1111 1111 1111 1111 1111 1100<sub>two</sub>





# Peer Instruction Question

---

$X = 1111\ 1111\ 1111\ 1111\ 1111\ 1111\ 1111\ 1100_{\text{two}}$

$Y = 0011\ 1011\ 1001\ 1010\ 1000\ 1010\ 0000\ 0000_{\text{two}}$

A.  $X > Y$  (if signed)

B.  $X > Y$  (if unsigned)

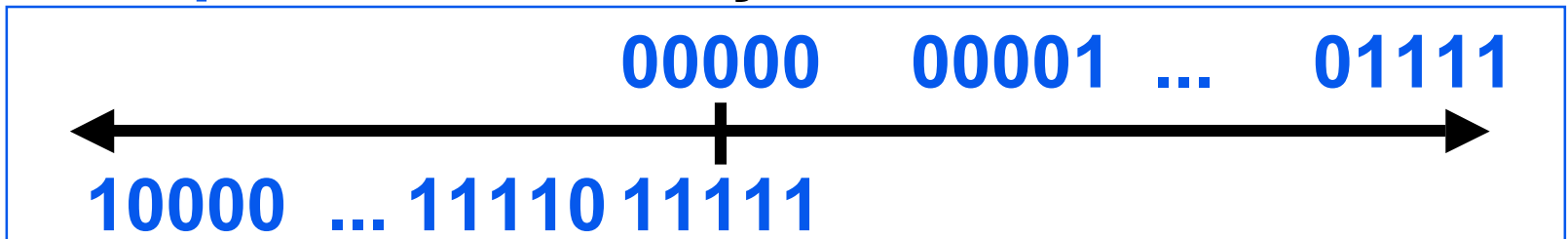
C. An encoding for Babylonians could have  $2^N$  non-zero numbers w/N bits!

	ABC
1:	FFF
2:	FFT
3:	FTF
4:	FTT
5:	TFF
6:	TFT
7:	TF
8:	TTT

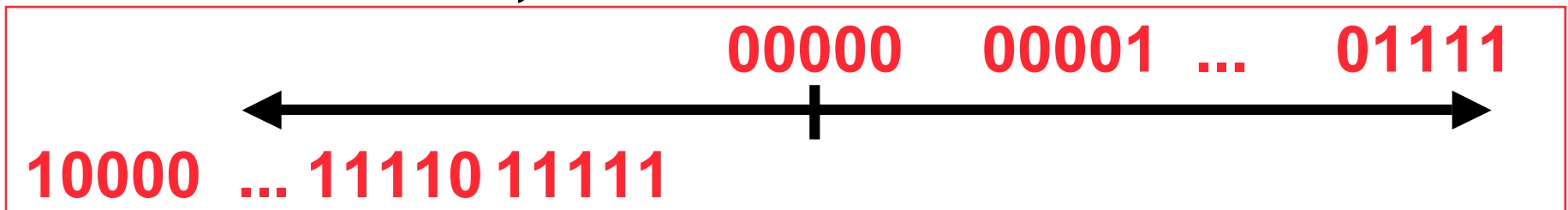


# And in Conclusion...

- We represent “things” in computers as particular bit patterns:  $N \text{ bits} \Rightarrow 2^N$
- Decimal for human calculations, binary for computers, hex to write binary more easily
- **1's complement** - mostly abandoned



- **2's complement** universal in computing: cannot avoid, so learn



• **Overflow: numbers  $\infty$ ; computers finite, errors!**

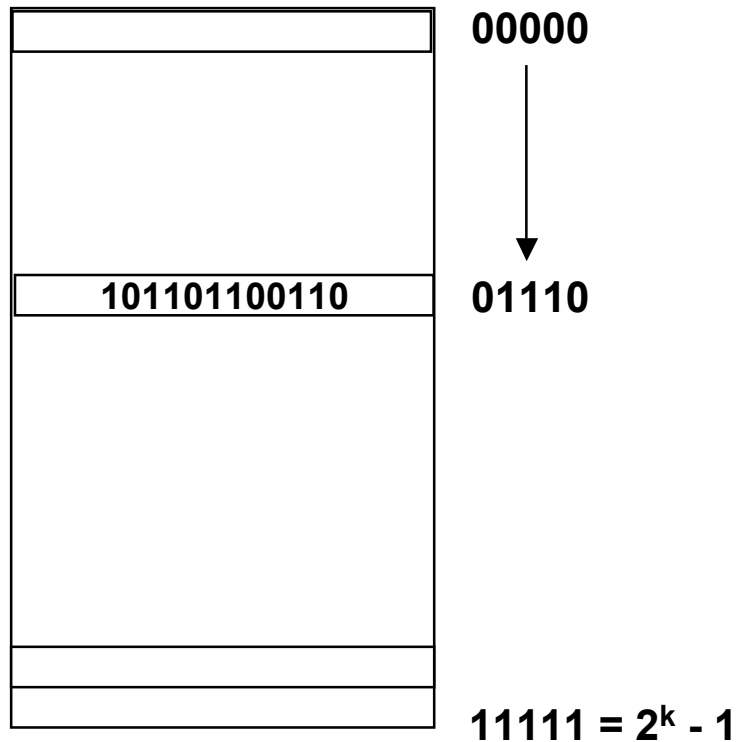
# Bonus Slides

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- **Peer instruction let's us skip example slides since you are expected to read book and lecture notes beforehand, but we include them for your review**
- **Slides shown in logical sequence order**



# BONUS: Numbers represented in memory



- **Memory is a place to store bits**
- **A *word* is a fixed number of bits (eg, 32) at an address**
- ***Addresses* are naturally represented as unsigned numbers in C**



## **BONUS: Signed vs. Unsigned Variables**

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- **Java just declares integers `int`**
  - **Uses two's complement**
- **C has declaration `int` also**
  - **Declares variable as a signed integer**
  - **Uses two's complement**
- **Also, C declaration `unsigned int`**
  - **Declares a unsigned integer**
  - **Treats 32-bit number as unsigned integer, so most significant bit is part of the number, not a sign bit**

