

Lecture 12 – Introduction to MIPS
Procedures II, Logical and Shift Ops



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Lecturer PSOE Dan Garcia

www.cs.berkeley.edu/~ddgarcia

Gotta love Sept/Oct! ⇒

Pennant races heating up!
SF, As trying to make the post-season.
Yankees trying to hold off Boston.
Ichiro soon to beat 84yr hit record!
espn.com/mlb/



Review

- Functions called with `jal`, return with `jr $ra`.
- The stack is your friend: Use it to save anything you need. Just be sure to leave it the way you found it.
- Instructions we know so far
 - Arithmetic: `add`, `addi`, `sub`, `addu`, `addiu`, `subu`
 - Memory: `lw`, `sw`
 - Decision: `beq`, `bne`, `slt`, `slti`, `sltu`, `sltiu`
 - Unconditional Branches (Jumps): `j`, `jal`, `jr`

- Registers we know so far

• All of them!

• There are CONVENTIONS when calling procedures!



Register Conventions (1/4)

- **Caller**: the calling function
- **Callee**: the function being called
- When callee returns from executing, the caller needs to know which registers may have changed and which are guaranteed to be unchanged.
- **Register Conventions**: A set of generally accepted rules as to which registers will be unchanged after a procedure call (`jal`) and which may be changed.



Register Conventions (2/4) - saved

- `$0`: **No Change**. Always 0.
- `$s0-$s7`: **Restore if you change**. Very important, that's why they're called saved registers. If the callee changes these in any way, it must restore the original values before returning.
- `$sp`: **Restore if you change**. The stack pointer must point to the same place before and after the `jal` call, or else the caller won't be able to restore values from the stack.
- **HINT** -- All saved registers start with **S**!



Register Conventions (3/4) - volatile

- `$ra`: **Can Change**. The `jal` call itself will change this register. Caller needs to save on stack if nested call.
- `$v0-$v1`: **Can Change**. These will contain the new returned values.
- `$a0-$a3`: **Can change**. These are volatile argument registers. Caller needs to save if they'll need them after the call.
- `$t0-$t9`: **Can change**. That's why they're called temporary; any procedure may change them at any time. Caller needs to save if they'll need them afterwards.



Register Conventions (4/4)

- What do these conventions mean?
 - If function R calls function E, then function R must save any temporary registers that it may be using onto the stack before making a `jal` call.
 - Function E must save any S (saved) registers it intends to use before garbling up their values
 - Remember: Caller/callee need to save only temporary/saved registers **they are using**, not all registers.



Parents leaving for weekend analogy (1/5)

- Parents (**main**) leaving for weekend
- They (**caller**) give keys to the house to kid (**callee**) with the rules (**calling conventions**):
 - You can trash the temporary room(s), like the den and basement (**registers**) if you want, we don't care about it
 - **BUT** you'd better leave the rooms (**registers**) that we want to **save** for the guests untouched. "these rooms better look the same when we return!"



Who hasn't heard this in their life?

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Parents leaving for weekend analogy (2/5)

- Kid now "owns" rooms (**registers**)
- Kid wants to use the **saved** rooms for a wild, wild party (**computation**)
- What does kid (**callee**) do?
 - Kid takes what was in these rooms and puts them in the garage (**memory**)
 - Kid throws the party, **trashes everything** (except garage, who goes there?)
 - Kid restores the rooms the parents wanted **saved after the party** by **replacing the items from the garage (**memory**) back into those saved rooms**



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Parents leaving for weekend analogy (3/5)

- Same scenario, except **before** parents return and kid replaces **saved** rooms...
- Kid (**callee**) has left valuable stuff (**data**) all over.
 - Kid's friend (**another callee**) wants the house for a party when the **kid** is away
 - Kid knows that friend might **trash the place** destroying valuable stuff!
 - Kid remembers rule parents taught and now becomes the "heavy" (**caller**), instructing friend (**callee**) on good rules (**conventions**) of house.



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Parents leaving for weekend analogy (4/5)

- If kid had data in **temporary rooms** (which were going to be trashed), there are three options:
 - Move items directly to garage (**memory**)
 - Move items to **saved rooms** whose contents have already been moved to the garage (**memory**)
 - Optimize lifestyle (**code**) so that the amount you've got to shlep stuff back and forth from garage (**memory**) is minimized
- Otherwise: "Dude, where's my data?!"



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Parents leaving for weekend analogy (5/5)

- **Friend** now "owns" rooms (**registers**)
- Friend wants to use the **saved** rooms for a wild, wild party (**computation**)
- What does friend (**callee**) do?
 - Friend takes what was in these rooms and puts them in the garage (**memory**)
 - Friend throws the party, **trashes everything** (except garage)
 - Friend restores the rooms the kid wanted **saved after the party** by **replacing the items from the garage (**memory**) back into those saved rooms**



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Bitwise Operations

- Up until now, we've done arithmetic (**add, sub, addi**), memory access (**lw and sw**), and branches and jumps.
- All of these instructions view contents of register as a single quantity (such as a signed or unsigned integer)
- **New Perspective:** View register as 32 raw bits rather than as a single 32-bit number
- Since registers are composed of 32 bits, we may want to access individual bits (or groups of bits) rather than the whole.
- Introduce two new classes of instructions:
 - Logical & Shift Ops



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Logical Operators (1/3)

- Two basic logical operators:
 - AND: outputs 1 only if **both** inputs are 1
 - OR: outputs 1 if **at least one** input is 1
- Truth Table: standard table listing all possible combinations of inputs and resultant output for each. E.g.,

A	B	A AND B	A OR B
0	0	0	0
0	1	0	1
1	0	0	1
1	1	1	1



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Logical Operators (2/3)

- Logical Instruction Syntax:
 - 1 2,3,4
 - where
 - 1) operation name
 - 2) register that will receive value
 - 3) first operand (register)
 - 4) second operand (register) or immediate (numerical constant)
- In general, can define them to accept > 2 inputs, but in the case of MIPS assembly, these accept exactly 2 inputs and produce 1 output



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- Again, rigid syntax, simpler hardware

Logical Operators (3/3)

- Instruction Names:
 - and, or: Both of these expect the third argument to be a register
 - andi, ori: Both of these expect the third argument to be an immediate
- MIPS Logical Operators are all **bitwise**, meaning that bit 0 of the output is produced by the respective bit 0's of the inputs, bit 1 by the bit 1's, etc.
 - C: Bitwise AND is & (e.g., $z = x \& y;$)
 - C: Bitwise OR is | (e.g., $z = x | y;$)



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Uses for Logical Operators (1/3)

- Note that anding a bit with 0 produces a 0 at the output while anding a bit with 1 produces the original bit.

- This can be used to create a **mask**.

- Example:

```
1011 0110 1010 0100 0011 1101 1001 1010
mask:0000 0000 0000 0000 0000 1111 1111 1111
• The result of anding these:
0000 0000 0000 0000 0000 1101 1001 1010
mask last 12 bits
```



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Uses for Logical Operators (2/3)

- The second bitstring in the example is called a **mask**. It is used to isolate the rightmost 12 bits of the first bitstring by masking out the rest of the string (e.g. setting it to all 0s).
- Thus, the and operator can be used to set certain portions of a bitstring to 0s, while leaving the rest alone.
 - In particular, if the first bitstring in the above example were in \$t0, then the following instruction would mask it:

```
andi $t0, $t0, 0xFFFF
```



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Uses for Logical Operators (3/3)

- Similarly, note that oring a bit with 1 produces a 1 at the output while oring a bit with 0 produces the original bit.

- This can be used to force certain bits of a string to 1s.

- For example, if \$t0 contains 0x12345678, then after this instruction:

```
ori $t0, $t0, 0xFFFF
```

- ... \$t0 contains 0x1234FFFF (e.g. the high-order 16 bits are untouched, while the low-order 16 bits are forced to 1s).



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Shift Instructions (1/4)

- Move (shift) all the bits in a word to the left or right by a number of bits.

- Example: shift right by 8 bits

0001 0010 0011 0100 0101 0110 0111 1000

0000 0000 0001 0010 0011 0100 0101 0110

- Example: shift left by 8 bits

0001 0010 0011 0100 0101 0110 0111 1000

0011 0100 0101 0110 0111 1000 0000 0000



Shift Instructions (2/4)

- Shift Instruction Syntax:

1 2,3,4

- where

- 1) operation name
- 2) register that will receive value
- 3) first operand (register)
- 4) shift amount (constant < 32)

- MIPS shift instructions:

1. **sll** (shift left logical): shifts left and **fills emptied bits with 0s**
2. **srl** (shift right logical): shifts right and **fills emptied bits with 0s**
3. **sra** (shift right arithmetic): shifts right and **fills emptied bits by sign extending**



Shift Instructions (3/4)

- Example: shift right arith by 8 bits

0001 0010 0011 0100 0101 0110 0111 1000

0000 0000 0001 0010 0011 0100 0101 0110

- Example: shift right arith by 8 bits

1001 0010 0011 0100 0101 0110 0111 1000

1111 1111 1001 0010 0011 0100 0101 0110



Shift Instructions (4/4)

- Since shifting may be faster than multiplication, a good compiler usually notices when C code multiplies by a power of 2 and compiles it to a shift instruction:

a *= 8; (in C)

would compile to:

sll \$s0,\$s0,3 (in MIPS)

- Likewise, shift right to divide by powers of 2

- remember to use sra



Peer Instruction

```
r: ... # R/W $s0,$v0,$t0,$a0,$sp,$ra,mem
...   ### PUSH REGISTER(S) TO STACK?
jal e # Call e
...   # R/W $s0,$v0,$t0,$a0,$sp,$ra,mem
jr $ra # Return to caller of r

e: ... # R/W $s0,$v0,$t0,$a0,$sp,$ra,mem
jr $ra # Return to r
```

What does r have to push on the stack before "jal e"?

- 1: Nothing
- 2: 1 of (\$s0, \$sp, \$v0, \$t0, \$a0, \$ra)
- 3: 2 of (\$s0, \$sp, \$v0, \$t0, \$a0, \$ra)
- 4: 3 of (\$s0, \$sp, \$v0, \$t0, \$a0, \$ra)
- 5: 4 of (\$s0, \$sp, \$v0, \$t0, \$a0, \$ra)
- 6: 5 of (\$s0, \$sp, \$v0, \$t0, \$a0, \$ra)



"And in Conclusion..."

- **Register Conventions:** Each register has a purpose and limits to its usage. Learn these and follow them, even if you're writing all the code yourself.

- Logical and Shift Instructions

- Operate on bits individually, unlike arithmetic, which operate on entire word.
- Use to isolate fields, either by masking or by shifting back and forth.
- Use **shift left logical**, sll, for multiplication by powers of 2
- Use **shift right arithmetic**, sra, for division by powers of 2.

- New Instructions:

and,andi, or,ori, sll,srl,sra

