

Lecture 39
I/O : Networks

2004-12-01



Lecturer PSOE Dan Garcia

www.cs.berkeley.edu/~ddgarcia

Clean to Zombie Bot in 4min

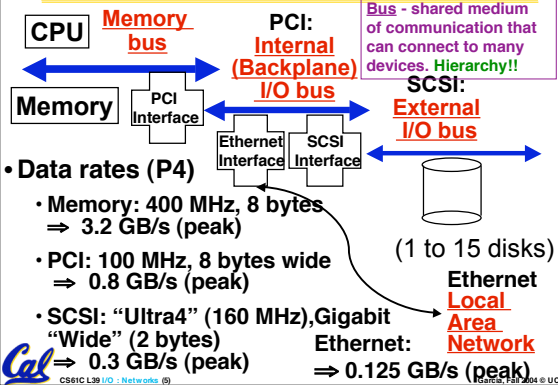
USA Today and Avantgarde report that it took less than 4 min for an unprotected PC running XP SP1 to be compromised. The Mac and Linux box were attacked but didn't fall.



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Buses in a PC: connect a few devices (2002)

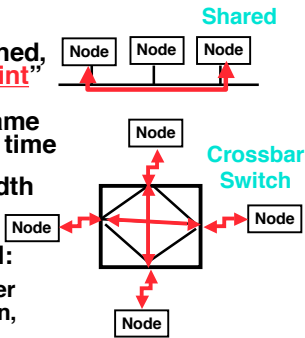


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Shared vs. Switched Based Networks

- **Shared Media vs. Switched:** in switched, pairs ("point-to-point" connections) communicate at same time; shared 1 at a time
- **Aggregate bandwidth (BW) in switched network is many times shared:**
 - point-to-point faster since no arbitration, simpler interface



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Why Networks?

- Originally sharing I/O devices between computers (e.g., printers)
- Then Communicating between computers (e.g. file transfer protocol)
- Then Communicating between people (e.g., email)
- Then Communicating between networks of computers ⇒ File sharing, WWW, ...



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How Big is the Network (1999)?

- ~30 Computers in 273 Soda
- ~400 in inst.cs.berkeley.edu
- ~4,000 in eeecs&cs .berkeley.edu
- ~50,000 in berkeley.edu
- ~5,000,000 in .edu
- ~46,000,000 in US (.com .net .edu .mil .us .org)
- ~56,000,000 in the world

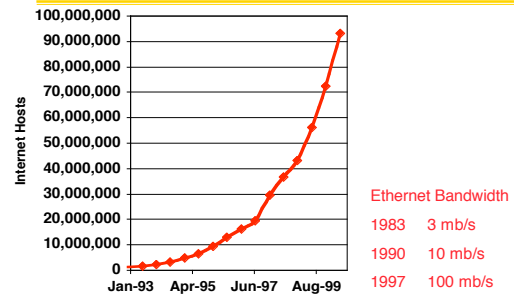


Source: Internet Software Consortium

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Growth Rate



"Source: Internet Software Consortium (<http://www.isc.org/>)".

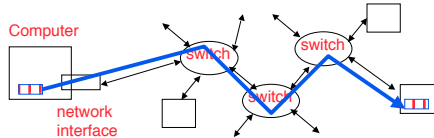


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What makes networks work?

- **links** connecting **switches** to each other and to computers or devices



- ability to **name** the components and to **route** packets of information - messages - from a source to a destination



- Layering, protocols, and encapsulation as means of **abstraction** (61C big idea)



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Typical Types of Networks

• Local Area Network (Ethernet)

- Inside a building: Up to 1 km
- (peak) Data Rate: 10 Mbits/sec, 100 Mbits/sec, 1000 Mbits/sec (1.25, 12.5, 125 MBytes/s)
- Run, installed by network administrators

• Wide Area Network

- Across a continent (10km to 10000 km)
- (peak) Data Rate: 1.5 Mb/s to 10000 Mb/s
- Run, installed by telecommunications companies (Sprint, UUNet[MCI], AT&T)



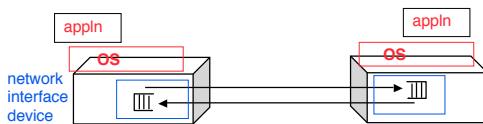
Wireless Networks (LAN), ...

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ABCs of Networks: 2 Computers

- **Starting Point**: Send bits between 2 computers



- Queue (First In First Out) on each end
- Can send both ways (“**Full Duplex**”)
- Information sent called a “**message**”
 - Note: Messages also called **packets**



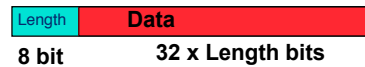
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A Simple Example: 2 Computers

• What is Message Format?

- Similar idea to Instruction Format
- Fixed size? Number bits?



- **Header(Trailer)**: information to deliver message
- **Payload**: data in message
- What can be in the data?
 - anything that you can represent as bits
 - values, chars, commands, addresses...



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Questions About Simple Example

• What if more than 2 computers want to communicate?

- Need computer “**address field**” in packet to know which computer should receive it (destination), and to which computer it came from for reply (source) [just like envelopes!]

Dest. Source Len

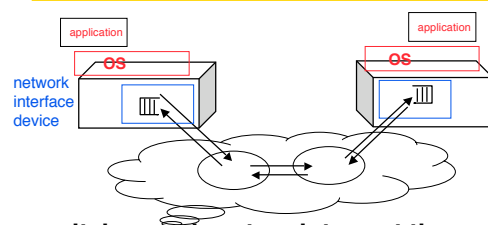
Net ID	Net ID	CMD/	Address	/Data
8 bits	8 bits	8 bits	32xn bits	
Header			Payload	



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ABCs: many computers



- switches and routers interpret the header in order to deliver the packet
- source encodes and destination decodes content of the payload



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Questions About Simple Example

- What if message is garbled in transit?
- Add redundant information that is checked when message arrives to be sure it is OK
- 8-bit sum of other bytes: called “**Check sum**”; upon arrival compare check sum to sum of rest of information in message. **xor** also popular.



Math 55 talks about what a Check sum is...

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Questions About Simple Example

- What if message never arrives?
- Receiver tells sender when it arrives (ack) [ala registered mail], sender retries if waits too long
- Don't discard message until get “ACK” (for ACKnowledgment); Also, if check sum fails, don't send ACK



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Observations About Simple Example

- Simple questions such as those above lead to more complex procedures to send/receive message and more complex message formats
- **Protocol**: algorithm for properly sending and receiving messages (packets)



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Software Protocol to Send and Receive

• SW Send steps

- 1: Application copies data to OS buffer
- 2: OS calculates checksum, starts timer
- 3: OS sends data to network interface HW and says start

• SW Receive steps

- 3: OS copies data from network interface HW to OS buffer
- 2: OS calculates checksum, if OK, send ACK; if not, **delete message** (sender resends when timer expires)
- 1: If OK, OS copies data to user address space, & signals application to continue



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Protocol for Networks of Networks?

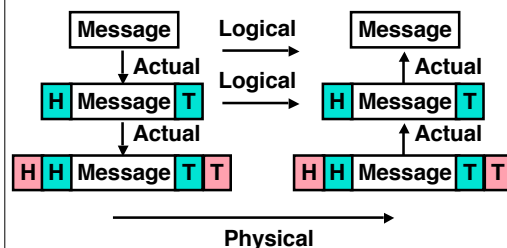
- **Internetworking**: allows computers on independent and incompatible networks to communicate reliably and efficiently;
 - Enabling technologies: SW standards that allow reliable communications without reliable networks
 - Hierarchy of SW layers, giving each layer responsibility for portion of overall communications task, called **protocol families** or **protocol suites**
- **Abstraction** to cope with **complexity of communication** vs. Abstraction for complexity of **computation**



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Protocol Family Concept



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Protocol Family Concept

- Key to **protocol families** is that communication occurs **logically** at the same level of the protocol, called **peer-to-peer**...

...but is **implemented via services at the next lower level**

- **Encapsulation**: carry higher level information within lower level “envelope”
- **Fragmentation**: break packet into multiple smaller packets and reassemble



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Protocol for Network of Networks

• **Transmission Control Protocol/Internet Protocol (TCP/IP)**

- This protocol family is the **basis of the Internet**, a WAN protocol
- IP makes best effort to deliver
- TCP guarantees delivery
- TCP/IP so popular it is used even when communicating locally: even across homogeneous LAN

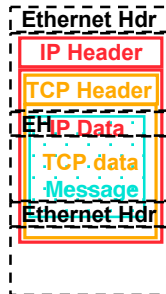


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TCP/IP packet, Ethernet packet, protocols

- Application sends message
- TCP breaks into 64KB segments, adds 20B header
- IP adds 20B header, sends to network
- If Ethernet, broken into 1500B packets with headers, trailers (24B)
- All Headers, trailers have length field, destination, ...



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Overhead vs. Bandwidth

- Networks are typically advertised using peak bandwidth of network link: e.g., 100 Mbits/sec Ethernet (“100 base T”)
- Software overhead to put message into network or get message out of network often limits useful bandwidth
- Assume overhead to send and receive = 320 microseconds (μ s), want to send 1000 Bytes over “100 Mbit/s” Ethernet

- Network transmission time:
 $1000\text{B} \times 8\text{b/B} / 100\text{Mb/s}$
 $= 8000\text{b} / (100\text{b}/\mu\text{s}) = 80 \mu\text{s}$

Effective bandwidth: $8000\text{b} / (320 + 80)\mu\text{s} = 20 \text{ Mb/s}$



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Peer Instruction

(T / F) P2P filesharing has been the dominant application on many links!

Suppose we have 2 networks, Which has a higher effective bandwidth as a function of the transferred data size?

- **BearsNet**
TCP/IP overhead 300 μ s, peak BW 10Mb/s
- **CalNet**
TCP/IP overhead 500 μ s, peak BW 100Mb/s

- | | |
|----|-----------|
| | TRUE |
| 1: | B always |
| 2: | C always |
| 3: | B small |
| | C big |
| 4: | B big |
| | C small |
| 5: | The same! |
| | FALSE |
| 6: | B always |
| 7: | C always |
| 8: | B small |
| | C big |
| 9: | B big |
| | C small |
| 0: | The same! |

And in conclusion...

- Protocol suites allow heterogeneous networking
 - Another form of principle of abstraction
 - Protocols \Rightarrow operation in presence of failures
 - Standardization key for LAN, WAN
- Integrated circuit (“Moore’s Law”) revolutionizing network switches as well as processors
 - Switch just a specialized computer
- Trend from shared to switched networks to get faster links and scalable bandwidth



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