inst.eecs.berkeley.edu/~cs61c

CS61C : Machine Structures

Lecture 40
Performance I

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Hybrid Hard Drives (HHT) ⇒ Samsung & MS announced new drives, which would use flash memory to cache information on disk, so the drive could spin down & save power when on, as well as boot much faster.

www.samsung.com/PressCenter/PressRelease/PressRelease.asp?seq=20050425_0000116210

Cool addition to the last lecture

- Drives inside the iPod and iPod Mini:

Thanks to Andy Dahl for the tip

Hitachi 1 inch 4GB MicroDrive

Toshiba 1.8-inch 20/40/60GB (MK1504GAL)

Review

- Magnetic disks continue rapid advance: 2x/yr capacity, 2x/2-yr bandwidth, slow on seek, rotation improvements, MB/$ 2x/yr!
  - Designs to fit high volume form factor
- RAID
  - Motivation: In the 1980s, there were 2 classes of drives: expensive, big for enterprises and small for PCs. They thought "make one big out of many small!"
  - Higher performance with more disk arms per $
  - Adds option for small # of extra disks (the "R")
  - Started @ Cal by CS Profs Katz & Patterson

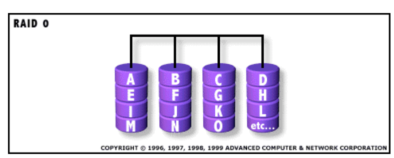
Redundant Arrays of (Inexpensive) Disks

- Files are "striped" across multiple disks
- Redundancy yields high data availability
  - Availability: service still provided to user, even if some components failed
- Disks will still fail
- Contents reconstructed from data redundantly stored in the array
  - ⇒ Capacity penalty to store redundant info
  - ⇒ Bandwidth penalty to update redundant info

Berkeley History, RAID-I

- RAID-I (1989)
  - Consisted of a Sun 4/280 workstation with 128 MB of DRAM, four dual-string SCSI controllers, 28 5.25-inch SCSI disks and specialized disk striping software
- Today RAID is > $27 billion dollar industry, 80% nonPC disks sold in RAIDs

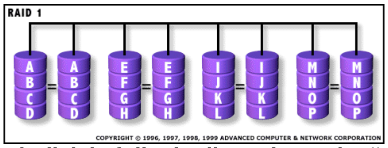
"RAID 0": No redundancy = "AID"



- Assume have 4 disks of data for this example, organized in blocks
- Large accesses faster since transfer from several disks at once

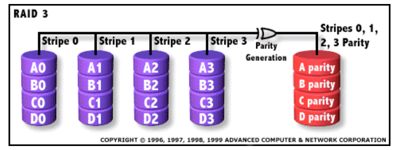
This and next 5 slides from RAID.edu, http://www.acnc.com/04_01_00.html

RAID 1: Mirror data



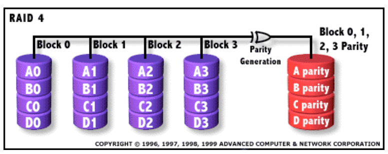
- Each disk is fully duplicated onto its "mirror"
  - Very high availability can be achieved
- Bandwidth reduced on write:
  - 1 Logical write = 2 physical writes
- Most expensive solution: 100% capacity overhead

RAID 3: Parity



- Parity computed across group to protect against hard disk failures, stored in P disk
- Logically, a single high capacity, high transfer rate disk
- 25% capacity cost for parity in this example vs. 100% for RAID 1 (5 disks vs. 8 disks)

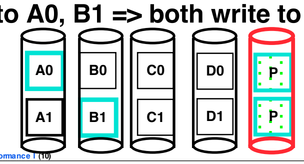
RAID 4: parity plus small sized accesses



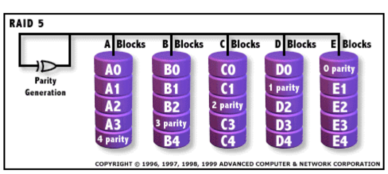
- RAID 3 relies on parity disk to discover errors on Read
- But every sector has an error detection field
- Rely on error detection field to catch errors on read, not on the parity disk
- Allows small independent reads to different disks simultaneously

Inspiration for RAID 5

- Small writes (write to one disk):
  - Option 1: read other data disks, create new sum and write to Parity Disk (access all disks)
  - Option 2: since P has old sum, compare old data to new data, add the difference to P: 1 logical write = 2 physical reads + 2 physical writes to 2 disks
- Parity Disk is bottleneck for Small writes: Write to A0, B1 => both write to P disk

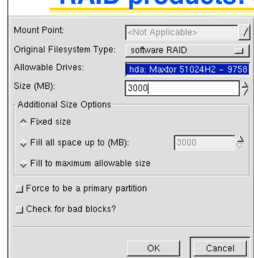


RAID 5: Rotated Parity, faster small writes



- Independent writes possible because of interleaved parity
  - Example: write to A0, B1 uses disks 0, 1, 4, 5, so can proceed in parallel
  - Still 1 small write = 4 physical disk accesses

RAID products: Software, Chips, Systems



RAID is $32 B industry in 2002, 80% nonPC disks sold in RAIDs

Margin of Safety in CS&E?

- Patterson reflects…
  - Operator removing good disk vs. bad disk
  - Temperature, vibration causing failure before repair
  - In retrospect, suggested RAID 5 for what we anticipated, but should have suggested RAID 6 (double failure OK) for unanticipated/safety margin…

Peer Instruction

1. RAID 1 (mirror) and 5 (rotated parity) help with performance and availability
2. RAID 1 has higher cost than RAID 5
3. Small writes on RAID 5 are slower than on RAID 1

	ABC
1:	FFF
2:	FFT
3:	FTF
4:	FTT
5:	TFF
6:	TFT
7:	TTF
8:	TTT

Administrivia

- Last semester's final + answers online soon
- HKN evaluations next Monday
- Final survey in lab this week
- Final exam review
  - Sunday, 2005-05-08 in the aft (location TBA)
- Final exam
  - Saturday, 2005-05-14 @ 12:30-3:30pm (loc TBA)
  - Same rules as Midterm, except you get 2 double-sided handwritten review sheets (1 from your midterm, 1 new one) + green sheet

Upcoming Calendar

Week #	Mon	Wed	Thu Lab	Fri
#15 This week	Perfor-mance I	Perfor-mance II	I/O Networks	TA Andy TBD
#16 Next Week Sun aft Review	LAST CLASS Summary Review & HKN Evals			FINAL EXAM SAT 05-14 @ 12:30pm

Performance

- Purchasing Perspective: given a collection of machines (or upgrade options), which has the
  - best performance ?
  - least cost ?
  - best performance / cost ?
- Computer Designer Perspective: faced with design options, which has the
  - best performance improvement ?
  - least cost ?
  - best performance / cost ?
- All require basis for comparison and metric for evaluation
- Solid metrics lead to solid progress!

Two Notions of "Performance"

Plane	DC to Paris	Top Speed	Passen-gers	Throughput (pmph)
Boeing 747	6.5 hours	610 mph	470	286,700
BAD/Sud Concorde	3 hours	1350 mph	132	178,200

- Which has higher performance?
  - Time to deliver 1 passenger?
  - Time to deliver 400 passengers?
- In a computer, time for 1 job called Response Time or Execution Time
- In a computer, jobs per day called Throughput or Bandwidth

Definitions

- Performance is in units of things per sec
  - bigger is better
- If we are primarily concerned with response time
  - $\text{performance}(x) = \dfrac{1}{\text{execution\_time}(x)}$

" F(ast) is *n* times faster than S(low) " means…

$$n = \frac{\text{performance}(F)}{\text{performance}(S)} = \frac{\text{execution\_time}(S)}{\text{execution\_time}(F)}$$

Example of Response Time v. Throughput

- Time of Concorde vs. Boeing 747?
  - Concord is 6.5 hours / 3 hours = 2.2 times faster
- Throughput of Boeing vs. Concorde?
  - Boeing 747: 286,700 pmph / 178,200 pmph = 1.6 times faster
- Boeing is 1.6 times ("60%") faster in terms of throughput
- Concord is 2.2 times ("120%") faster in terms of flying time (response time)

We will focus primarily on execution time for a single job

Confusing Wording on Performance

- Will (try to) stick to "n times faster"; its less confusing than "m % faster"
- As faster means both increased performance and decreased execution time, to reduce confusion we will (and you should) use "improve performance" or "improve execution time"

What is Time?

- Straightforward definition of time:
  - Total time to complete a task, including disk accesses, memory accesses, I/O activities, operating system overhead, ...
  - "real time", "response time" or "elapsed time"
- Alternative: just time processor (CPU) is working only on your program (since multiple processes running at same time)
  - "CPU execution time" or "CPU time"
  - Often divided into system CPU time (in OS) and user CPU time (in user program)

How to Measure Time?

- User Time ⇒ seconds
- CPU Time: Computers constructed using a clock that runs at a constant rate and determines when events take place in the hardware
  - These discrete time intervals called clock cycles (or informally clocks or cycles)
  - Length of clock period: clock cycle time (e.g., 2 nanoseconds or 2 ns) and clock rate (e.g., 500 megahertz, or 500 MHz), which is the inverse of the clock period; use these!

"And in conclusion…"

- RAID
  - Motivation: In the 1980s, there were 2 classes of drives: expensive, big for enterprises and small for PCs. They thought "make one big out of many small!"
  - Higher performance with more disk arms per $
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  - Started @ Cal by CS Profs Katz & Patterson
- Latency v. Throughput
- Measure time as User time vs CPU time