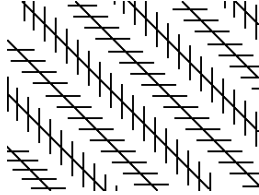


## Lecture #28: Parallel Computing



2005-08-09

Andy Carle



## Scientific Computing

- **Traditional Science**
  - 1) Produce theories and designs on “paper”
  - 2) Perform experiments or build systems
  - Has become difficult, expensive, slow, and dangerous for fields on the leading edge
- **Computational Science**
  - Use ultra-high performance computers to simulate the system we’re interested in
- **Acknowledgement**
  - Many of the concepts and some of the content of this lecture were drawn from Prof. Jim Demmel’s CS 267 lecture slides which can be found at [http://www.cs.berkeley.edu/~demmel/cs267\\_Spr05/](http://www.cs.berkeley.edu/~demmel/cs267_Spr05/)



## Example Applications

- **Science**
  - Global climate modeling
  - Biology: genomics; protein folding; drug design
  - Astrophysical modeling
  - Computational Chemistry
  - Computational Material Sciences and Nanosciences
- **Engineering**
  - Semiconductor design
  - Earthquake and structural modeling
  - Computation fluid dynamics (airplane design)
  - Combustion (engine design)
  - Crash simulation
- **Business**
  - Financial and economic modeling
  - Transaction processing, web services and search engines
- **Defense**
  - Nuclear weapons – test by simulations
  - Cryptography



## Performance Requirements

- **Terminology**
  - Flop – Floating point operation
  - Flops/second – standard metric for expressing the computing power of a system
- **Global Climate Modeling**
  - Divide the world into a grid (e.g. 10 km spacing)
  - Solve fluid dynamics equations to determine what the air has done at that point every minute
    - Requires about 100 Flops per grid point per minute
  - This is an extremely simplified view of how the atmosphere works, to be maximally effective you need to simulate many additional systems on a much finer grid



## Performance Requirements (2)

- **Computational Requirements**
  - To keep up with real time (i.e. simulate one minute per wall clock minute): 8 Gflops/sec
  - Weather Prediction (7 days in 24 hours): 56 Gflops/sec
  - Climate Prediction (50 years in 30 days): 4.8 Tflops/sec
  - Climate Prediction Experimentation (50 years in 12 hours): 288 Tflops/sec
- **Perspective**
  - Pentium 4 1.4GHz, 1GB RAM, 4x100MHz FSB
    - ~320 Mflops/sec, effective
    - Climate Prediction would take ~1233 years



## What Can We Do?

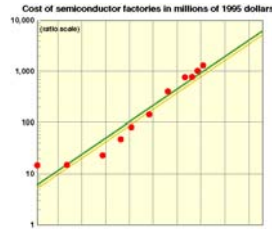
- **Wait**
  - Moore’s law tells us things are getting better; why not stall for the moment?
- **Parallel Computing!**



## Prohibitive Costs

### Rock's Law

- The cost of building a semiconductor chip fabrication plant that is capable of producing chips in line with Moore's law doubles every four years



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## How fast can a serial computer be?

- Consider a 1 Tflop/sec sequential machine:
  - Data must travel some distance,  $r$ , to get from memory to CPU
  - To get 1 data element per cycle, this means  $10^{12}$  times per second at the speed of light,  $c = 3 \times 10^8$  m/s. Thus  $r < c/10^{12} = 0.3$  mm
    - So all of the data we want to process must be stored within 0.3 mm of the CPU
- Now put 1 Tbyte of storage in a 0.3 mm x 0.3 mm area:
  - Each word occupies about 3 square Angstroms, the size of a very small atom
  - Maybe someday, but it most certainly isn't going to involve transistors as we know them



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## What is Parallel Computing?

- Dividing a task among multiple processors to arrive at a unified (meaningful) solution
  - For today, we will focus on systems with many processors executing identical code
- How is this different from Multiprogramming (which we've touched on some in this course)?
- How is this different from Distributed Computing?



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## Recent History

- Parallel Computing as a field exploded in popularity in the mid-1990s
- This resulted in an "arms race" between universities, research labs, and governments to have the fastest supercomputer in the world



Source: top500.org



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## Current Champions



BlueGene/L – IBM/DOE  
Rochester, United States  
32768 Processors, 70.72 Tflops/sec  
0.7 GHz PowerPC 440



Columbia – NASA/Ames  
Mountain View, United States  
10160 Processors, 51.87 Tflops/sec  
1.5 GHz SGI Altix



Earth Simulator – Earth Simulator Ctr.  
Yokohama, Japan  
5120 Processors, 35.86 Tflops/sec  
SX6 Vector



Data Source: top500.org  
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## Administrivia

- Proj 4 Due Friday
- HW8 (Optional) Due Friday
- Final Exam on Friday
  - Yeah, sure, you can have 3 one-sided cheat sheets
    - But I really don't think they'll help you all that much
- Course Survey in lab today



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## Parallel Programming

- Processes and Synchronization
- Processor Layout
- Other Challenges
  - Locality
  - Finding parallelism
  - Parallel Overhead
  - Load Balance



## Processes

- We need a mechanism to intelligently split the execution of a program

- Fork:

```
int main(...){
    int pid = fork();
    if (pid == 0) printf("I am the child.");
    if (pid != 0) printf("I am the parent.");
    return 0;
}
```

- What will this print?



## Processes (2)

- We don't know! Two potential orderings:
  - I am the child. I am the parent.
  - I am the parent. I am the child.
  - This situation is a simple race condition. This type of problem can get far more complicated...
- Modern parallel compilers and runtime environments hide the details of actually calling `fork()` and moving the processes to individual processors, but the complexity of synchronization remains



## Synchronization

- How do processors communicate with each other?
- How do processors know when to communicate with each other?
- How do processors know which other processor has the information they need?
- When you are done computing, which processor, or processors, have the answer?



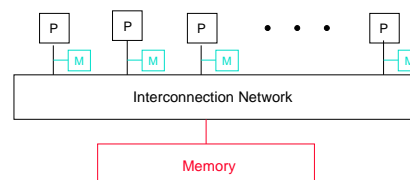
## Synchronization (2)

- Some of the logistical complexity of these operations is reduced by standard communication frameworks
  - Message Passing Interface (MPI)
- Sorting out the issue of who holds what data can be made easier with the use of explicitly parallel languages
  - Unified Parallel C (UPC)
  - Titanium (Parallel Java Variant)
- Even with these tools, much of the skill and challenge of parallel programming is in resolving these problems



## Processor Layout

### Generalized View

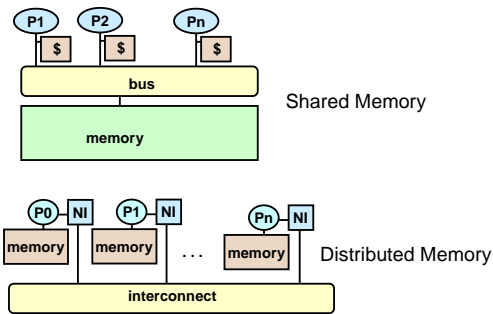


M = Memory local to one processor

Memory = Memory local to all other processors



## Processor Layout (2)



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## Processor Layout (3)

- Clusters of SMPs
  - n of the N total processors share one memory
  - Simple shared memory communication within one cluster of n processors
  - Explicit network type calls to communicate from one group of n to another
- Understanding the processor layout that your application will be running on is crucial!

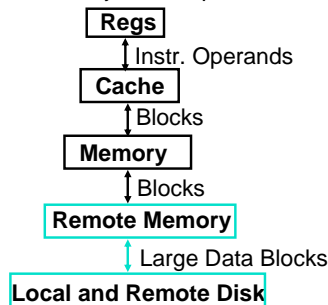


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## Parallel Locality

- We now have to expand our view of the memory hierarchy to include remote machines
- Remote memory behaves like a very fast network
  - Bandwidth vs. Latency becomes important



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## Amdahl's Law

- Applications can almost never be completely parallelized
  - Let s be the fraction of work done sequentially, so (1-s) is fraction parallelizable, and P = number of processors
- $$\text{Speedup}(P) = \text{Time}(1) / \text{Time}(P)$$
- $$\leq 1 / (s + (1-s)/P)$$
- $$\leq 1/s$$
- Even if the parallel portion of your application speeds up perfectly, your performance may be limited by the sequential portion



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## Parallel Overhead

- Given enough parallel work, these are the biggest barriers to getting desired speedup
- Parallelism overheads include:
  - cost of starting a thread or process
  - cost of communicating shared data
  - cost of synchronizing
  - extra (redundant) computation
- Each of these can be in the range of milliseconds (many millions of flops) on some systems
- Tradeoff: Algorithm needs sufficiently large units of work to run fast in parallel (i.e. large granularity), but not so large that there is not enough parallel work



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## Load Balance

- Load imbalance is the time that some processors in the system are idle due to
  - insufficient parallelism (during that phase)
  - unequal size tasks
- Examples of the latter
  - adapting to "interesting parts of a domain"
  - tree-structured computations
  - fundamentally unstructured problems
- Algorithms need to carefully balance load



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## Summary

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- **Parallel Computing is a multi-billion dollar industry driven by interesting and useful scientific computing applications**
- **It is extremely unlikely that sequential computing will ever again catch up with the processing power of parallel systems**
- **Programming parallel systems can be extremely challenging, but is built upon many of the concepts you've learned this semester in 61c**

